

# Improving the Dynamic Programming Algorithm for Nurse Rostering

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## Overview

### Existing algorithm

- A dynamic programming algorithm for optimally timetabling one nurse
- Found in column generation algorithms for nurse rostering
- Runs in time polynomial in the number of days to be timetabled

### This paper's contribution

- Generalize to multiple nurses, any subset of the days, all constraints
- Speed up (essential for multiple nurses)

Aim is to use it as reconstruction operator in VLSN search (still to do).

## **Nurse rostering**

Assign shifts (e.g. morning, afternoon, night) to the nurses of a hospital ward, over several weeks.

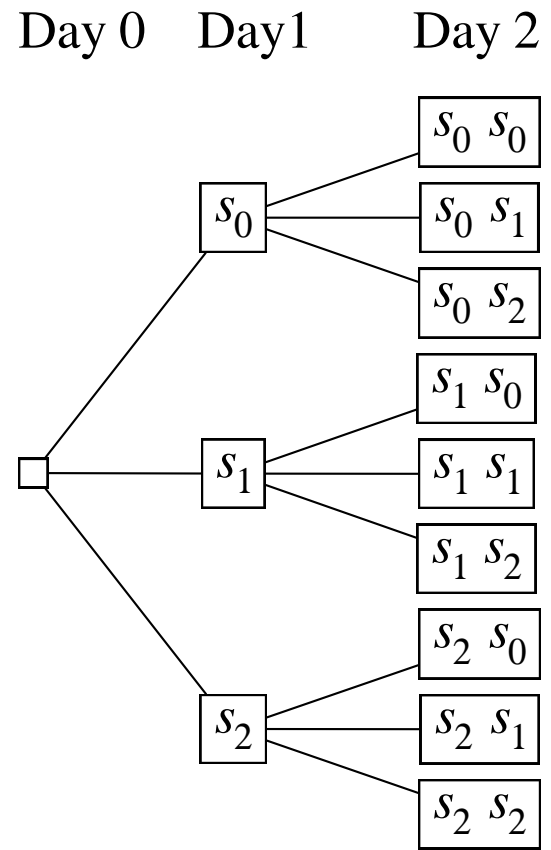
### **Cover constraints**

- For each shift, limits on the number of nurses assigned
- Requests for nurses may specify skills (senior, trainee, etc.)

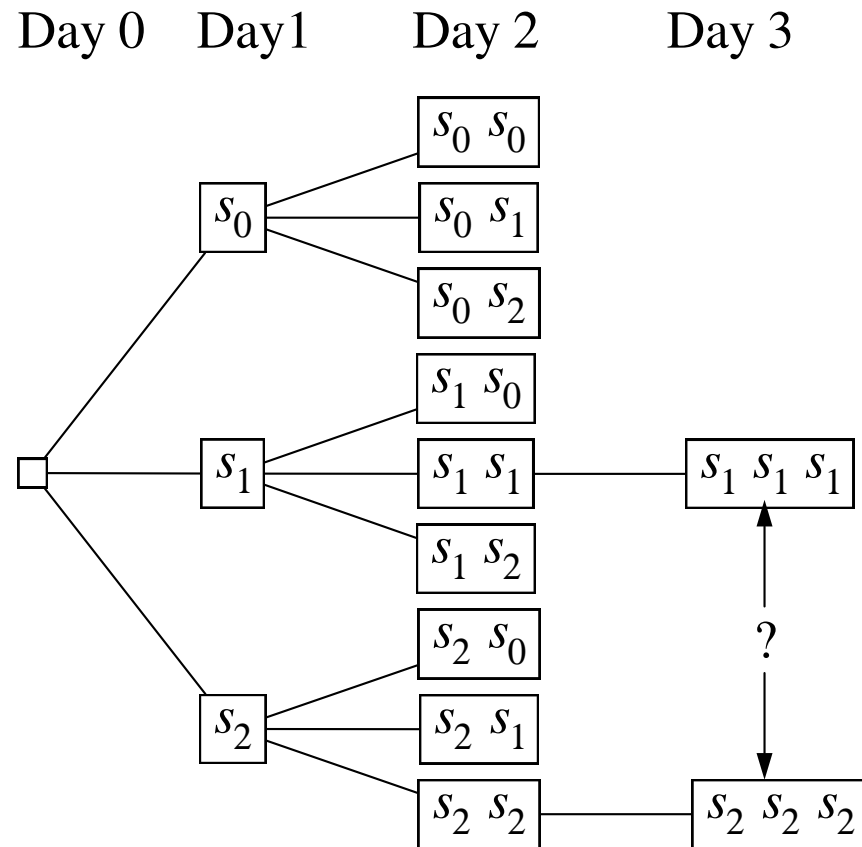
### **Resource constraints**

- Each nurse takes at most one shift per day (hard)
- Counter constraints: limits on total shifts, busy weekends, etc.
- Sequence constraints: limits on consecutive busy days, night shifts, etc.

### Simple tree search (single nurse)



## The dynamic programming algorithm



## Dominance testing

### Complete extension of solution $S$

A solution that begins with  $S$  and carries on to the end.

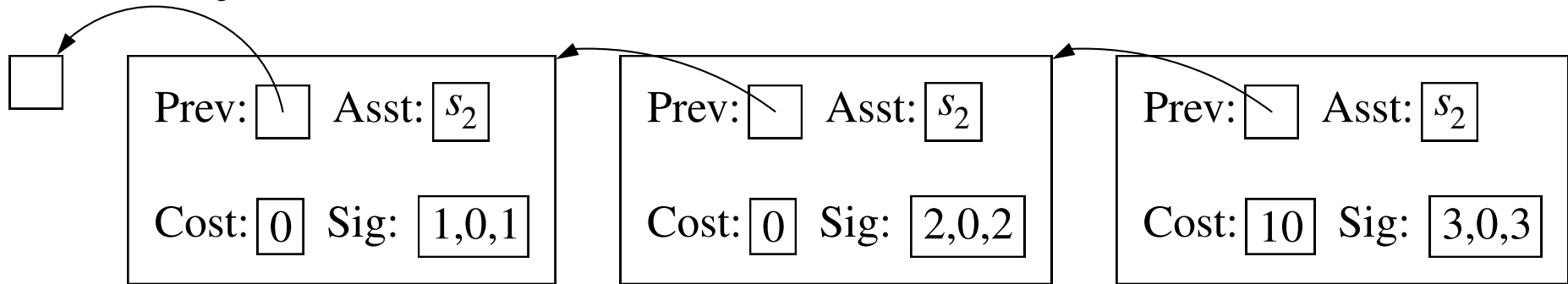
### Dominates( $S_1, S_2$ )

True when for each complete extension of  $S_2$  there is a complete extension of  $S_1$  of equal or less cost.

If Dominates( $S_1, S_2$ ) then we can delete  $S_2$ . Check this as each solution is created.

## How solutions are represented

### Solution objects (nodes of search tree)



### Signatures

Constraint	$s_1 \ s_1 \ s_1$	$s_2 \ s_2 \ s_2$
At most 5 shifts	3	3
At most 3 consecutive $s_1$ shifts	3	0
At most 2 consecutive $s_2$ shifts	0	3 + cost 10 (say)

## Basic dominance

$S_1$  dominates  $S_2$  when

- $cost(S_1) \leq cost(S_2)$  and
- for each constraint  $m$  in the signatures,
  - if upper limit only,  $sig(m, S_1) \leq sig(m, S_2)$
  - if lower limit only,  $sig(m, S_1) \geq sig(m, S_2)$
  - if both,  $sig(m, S_1) = sig(m, S_2)$

Seems to be the standard in the literature.



## Tradeoff dominance

- Seems to be new to this paper
- Similar to basic dominance
- But can trade off a small violation against  $cost(S_2) - cost(S_1)$

## Uniform dominance

- Done since paper was written
- Similar to tradeoff dominance
- Handles awkward cases well (allow zero flag, quadratic cost functions)
- May be best possible dominance test in practice

## Generalizing

### Multiple nurses

One assignment for each nurse, and all nurses' constraints in the signature:

Prev:	<input type="checkbox"/>	Asst:	$r_1 := s_2, r_2 := s_0$
Cost:	0	Sig:	2,0,2,2,2,0

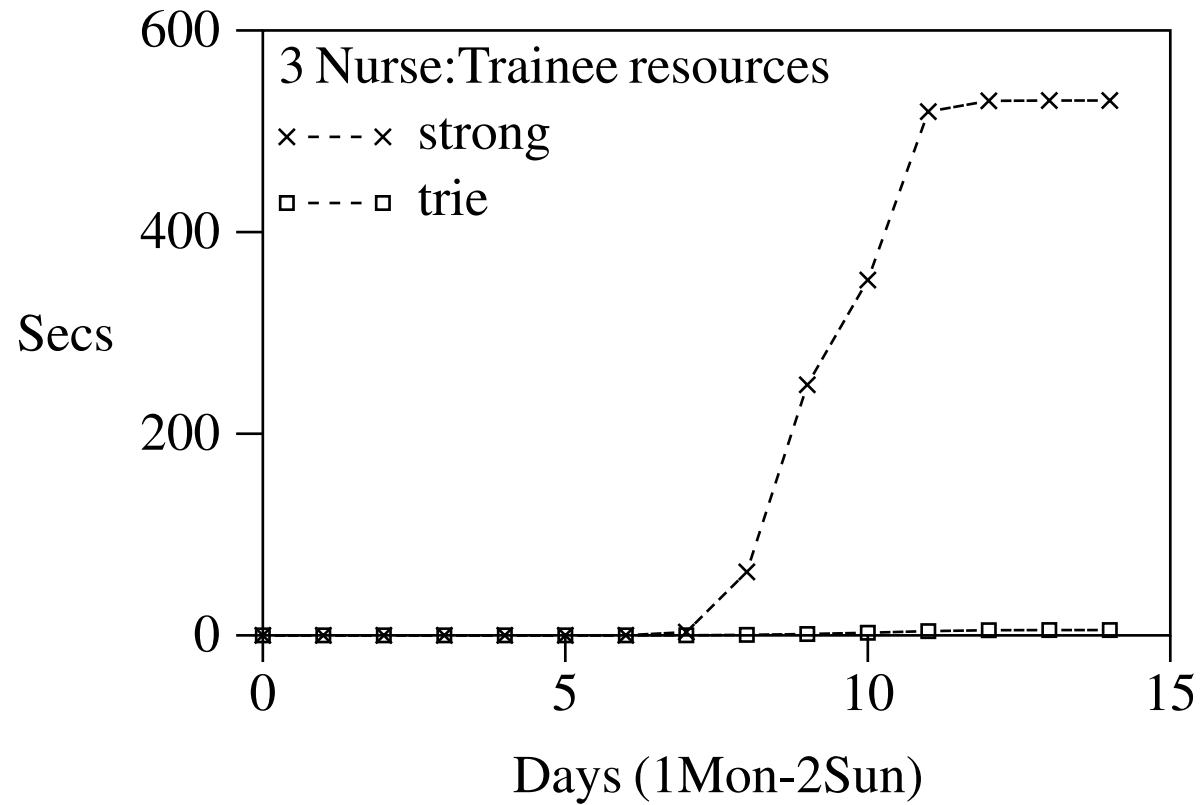
### Any subset of the days

Search over selected days only; unselected days add constants to signatures.

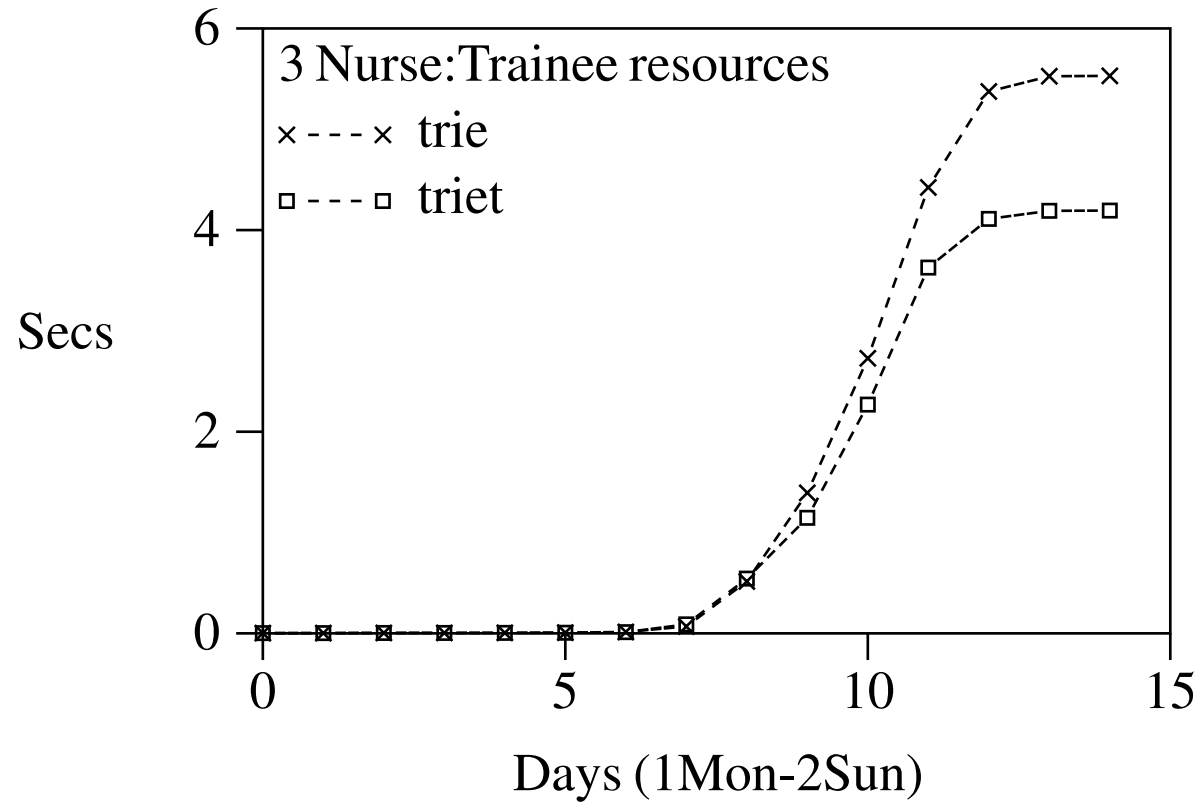
### All constraints

Support the XESTT format; its constraints cover all the well-known data sets.

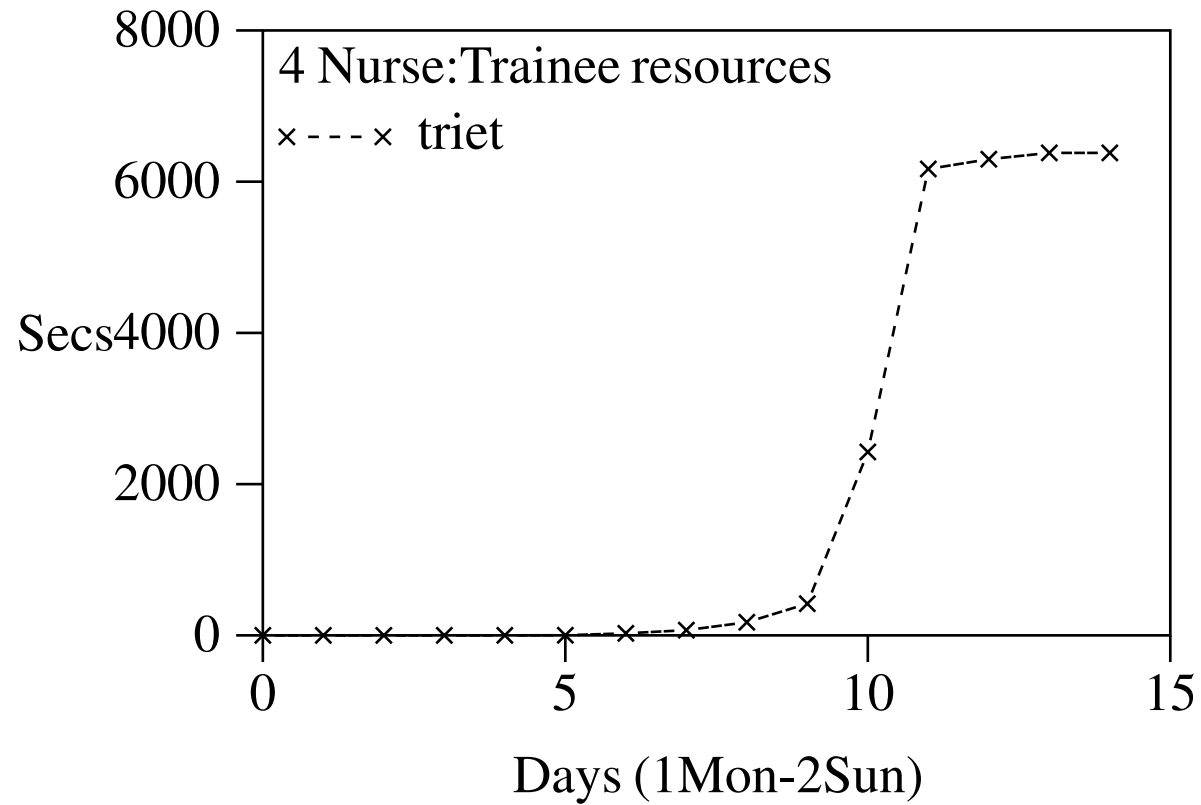
## Experiment 1 – Storing solutions in a trie data structure



## Experiment 2 – Moving to tradeoff dominance



### Experiment 3 – Four nurses not yet efficient



## **Conclusion**

### **What has been done**

- Algorithm generalized to arbitrary nurses, days, constraints
- And made to run much faster

### **What still needs to be done**

- More speedup (5 nurses at least)
- Test algorithm as VLSN reconstruction operator